Introduction

Syntax is constrained by LF, PF, and the parser — but is it also uniquely determined by them, or does it have some wiggle room as to how it satisfies their demands? For the parser, it might.

Reference-set computation (transderivationality) comes with a high resource load that fails the computability requirements of the parser. Thus it cannot be part of syntax. However, mathematics suggests that certain kinds of reference-set computation can be hidden from the parser.

A Mathematical Model of Reference-Set Computation

Reference-set constraints can be modelled as tree automata with outputs, which may be specified by Optimality Systems (a variant of OT grammars; Frank and Satta 1998).

An Optimality System consists of

- a collection of reference types over the input language,
- a collection of reference sets over the candidate language,
- a function mapping reference types to reference sets,
- a sequence of OT constraints.

Doing Away with Reference-Set Computation

Given our new perspective on reference-set computation, it is easy to show that some reference-set constraints can be reduced to standard well-formedness conditions (implemented as, say, additional features on lexical items etc.). This allows syntax to use any reference-set constraint for which there is an efficiently computable equivalent for the parser.

A reference-set constraint is reducible if the following holds for its respective Optimality System.

- Output joint preservation: If two reference sets overlap, then the reference types that are mapped to them overlap, too.
- Rigidity: If reference type \( T \) is mapped to reference-set \( R \), then an output candidate in \( R \) is optimal for some input of type \( T \) only if it is optimal for all inputs of type \( T \).
- The OT generator and each OT constraint can be modelled by finite-state tree automata with outputs.

Results for Constraints from the Literature

All prominent instances of reference-set computation (Fewest Steps, Rule I, Scope Economy, Stress Shift in Szendrői 2001) obey both output joint preservation and rigidity.

At least for Fewest Steps and Stress Shift, the conditions on the generator and the OT constraints are satisfied, too, so both constraints have efficiently computable equivalents that do not involve reference-set computation.

Conclusion and Open Questions

The reducibility of certain reference-set constraints to well-formedness conditions loosens the parser’s grip on syntax. The amount of reference-set computation in syntax depends only on the availability of computable equivalents.

This raises two intriguing questions:

- Are there similar means of escaping the interfaces’ clutches?
- Why should syntax prefer reference-set constraints over their efficiently computable doubles?

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