Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion

Morphotactics as Tier-Based Strictly Local Dependencies

Alëna Aksënova, Thomas Graf, and Sedigheh Moradi

Stony Brook University

SIGMORPHON 14 Berlin, Germany 11. August 2016

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Our goal				

Received view

Recent research

Phonology

regular Kaplan&Kay (1994) subregular Heinz (2015) Morphology regular Beesley&Karttunen (2003) ?

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Our goal				

	Phonology	Morphology
Received view	regular	regular
Received view	Kaplan&Kay (1994)	Beesley&Karttunen (2003)
Recent research	subregular Heinz (2015)	?

- Show that morphotactics is subregular
- More precisely: Tier-Based Strictly Local
- Consequences
 - parallels to phonology
 - learnable in the limit from positive text
 - explain typological gaps

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Outline				



- 2 SL Patterns In Morphology
- 3 Tier-Based Strictly Local
 - TSL is necessary
 - TSL is sufficient



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Morphotactics						

Definition (Morphotactics)

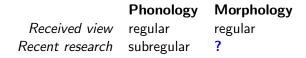
Restrictions on the linear ordering of morphemes.

- Our focus: morphotactics in underlying representations (English)
 OKSTEM-PL
 *PL-STEM
- \bullet \Rightarrow allomorphy (dogs, peaches) is not considered yet

 Preliminaries
 Lower bound (SL)
 Upper bound (TSL)
 Typology
 Conclusion

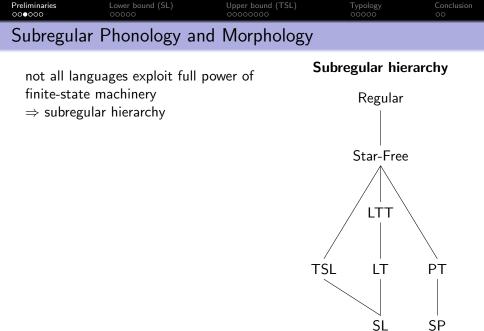
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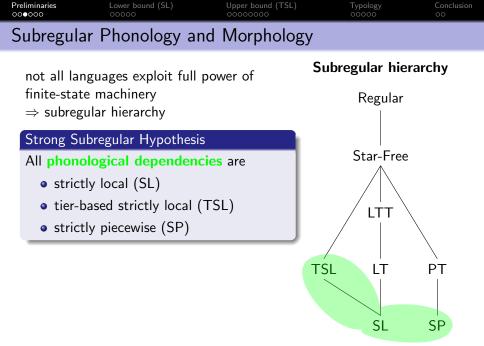
Computational nature of morphotactics

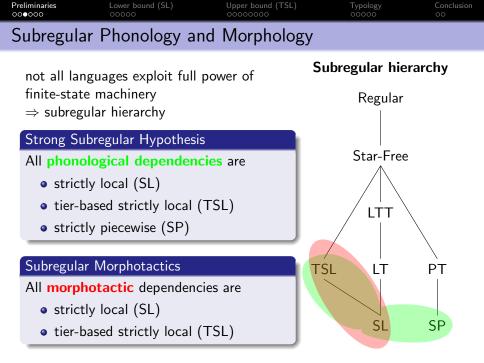


Advantages of (some) subregular languages:

- resolves learnability issues
- describes potential cognitive mechanisms
- uses less powerful generating device







Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Strictly L	ocal languages			

- SL and TSL are generated by *k*-gram models.
- A k-gram model is a finite set of blocked k-grams.

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Strictly L	ocal languages			

- SL and TSL are generated by *k*-gram models.
- A k-gram model is a finite set of blocked k-grams.

Example (Strictly Local Grammar for $(ab)^*a$)

 $\Sigma = \{a, b\}$ Grammar = {×b, bb, aa, b×, ××}

Accepted strings: $\forall a \ltimes, \forall a b a \ltimes, \forall a b a b a \ltimes, etc.$ Rejected strings: $\forall a b \ltimes, \forall b a \ltimes, \forall a b b a \ltimes, etc.$

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Accepted strings: $\forall a \ltimes, \forall a b a \ltimes, \forall a b a b a \ltimes, etc.$ Rejected strings: $\forall a \flat \ltimes, \forall b a \ltimes, \forall a b b a \ltimes, etc.$

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Accepted strings: $\forall a \ltimes, \forall a b a \ltimes, \forall a b a b a \ltimes, etc.$ Rejected strings: $\forall a \triangleright \ltimes, \forall b a \ltimes, \forall a b b a \ltimes, etc.$

Definition

Strictly *k*-**Local (***k*-**SL) grammar** consists of a set of blocked *k*-grams over an alphabet Σ .



Example (Tier-Based Strictly Local Grammar for $c^*(ac^*bc^*)^*ac^*$)

$$\begin{split} \Sigma &= \{ \mathsf{a}, \, \mathsf{b}, \, \mathsf{c} \} \\ \mathsf{Grammar:} \\ \mathsf{G}(\mathsf{a},\mathsf{b_tier}) &= \{ \rtimes \mathsf{b}, \, \mathsf{bb}, \, \mathsf{aa}, \, \mathsf{b} \ltimes, \, \rtimes \ltimes \} \end{split}$$

Accepted strings: $\forall a \ltimes, \forall accba \ltimes, \forall cacbaccccba \ltimes, etc.$



Example (Tier-Based Strictly Local Grammar for $c^*(ac^*bc^*)^*ac^*$)

$$\begin{split} \Sigma &= \{ \mathsf{a}, \, \mathsf{b}, \, \mathsf{c} \} \\ \mathsf{Grammar:} \\ \mathsf{G}(\mathsf{a}, \mathsf{b_tier}) &= \{ \rtimes \mathsf{b}, \, \mathsf{bb}, \, \mathsf{aa}, \, \mathsf{b} \ltimes, \, \rtimes \ltimes \} \\ \mathsf{Accepted strings:} \ &\rtimes \mathsf{a} \ltimes, \, \rtimes \mathsf{accba} \ltimes, \, \rtimes \mathsf{cacbaccccba} \ltimes, \, \mathsf{etc.} \\ & \mathsf{a}, \mathsf{b_tier:} \ &\rtimes \mathsf{a} \ltimes \ &\rtimes \mathsf{aba} \ltimes \ &\rtimes \mathsf{ababa} \ltimes \end{split}$$



Example (Tier-Based Strictly Local Grammar for $c^*(ac^*bc^*)^*ac^*$) $\Sigma = \{a, b, c\}$

Grammar:

 $\mathsf{G}(\mathsf{a},\mathsf{b_tier}) = \{ \rtimes \mathsf{b}, \ \mathsf{bb}, \ \mathsf{aa}, \ \mathsf{b\ltimes}, \ \rtimes \ltimes \}$

Accepted strings: $\forall a \ltimes$, $\forall accba \ltimes$, $\forall cacbaccccba \ltimes$, etc.

a,b_tier: ⋊*a*⋉ ⋊*aba*⋉ ⋊*ababa*⋉

Rejected strings: $\rtimes accccaba \ltimes$, $\rtimes abcccacccbc \ltimes$, etc.



Example (Tier-Based Strictly Local Grammar for $c^*(ac^*bc^*)^*ac^*$)

$$\begin{split} \Sigma &= \{ \texttt{a}, \texttt{b}, \texttt{c} \} \\ \texttt{Grammar:} \\ \texttt{G}(\texttt{a},\texttt{b_tier}) &= \{ \rtimes \texttt{b}, \texttt{bb}, \texttt{aa}, \texttt{b} \ltimes, \rtimes \ltimes \} \\ \texttt{Accepted strings:} &\rtimes \texttt{a} \ltimes, \rtimes \texttt{accba} \ltimes, \rtimes \texttt{cacbaccccba} \ltimes, \texttt{etc.} \\ & \texttt{a},\texttt{b_tier:} &\rtimes \texttt{a} \ltimes & \rtimes \texttt{abab} \ltimes & \rtimes \texttt{ababa} \ltimes \\ \texttt{Rejected strings:} &\rtimes \texttt{accccaba} \ltimes, \rtimes \texttt{abcccacccbc} \ltimes, \texttt{etc.} \\ & \texttt{a},\texttt{b_tier:} &\rtimes \texttt{aba} \ltimes & \rtimes \texttt{abab} \ltimes \\ \end{split}$$



Tier-Based Strictly Local languages

Example (Tier-Based Strictly Local Grammar for $c^*(ac^*bc^*)^*ac^*$)

$$\begin{split} \Sigma &= \{ \texttt{a}, \texttt{b}, \texttt{c} \} \\ \texttt{Grammar:} \\ \texttt{G}(\texttt{a},\texttt{b_tier}) &= \{ \rtimes \texttt{b}, \texttt{bb}, \texttt{aa}, \texttt{b} \ltimes, \rtimes \ltimes \} \\ \texttt{Accepted strings:} &\rtimes \texttt{a} \ltimes, \rtimes \texttt{accba} \ltimes, \rtimes \texttt{cacbaccccba} \ltimes, \texttt{etc.} \\ &\texttt{a},\texttt{b_tier:} &\rtimes \texttt{a} \ltimes &\rtimes \texttt{aba} \ltimes &\rtimes \texttt{ababa} \ltimes \\ \texttt{Rejected strings:} &\rtimes \texttt{accccaba} \ltimes, \rtimes \texttt{abcccacccbc} \ltimes, \texttt{etc.} \\ &\texttt{a},\texttt{b_tier:} &\rtimes \texttt{aaba} \ltimes &\rtimes \texttt{abab} \ltimes \end{split}$$

Definition

A Tier-Based Strictly k-Local grammar is a k-SL grammar that operates over a *tier*, a specific substructure of the string.

Learnability				
Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
00000●		00000000	00000	00

Learning of SL and TSL

- learning ≡ memorizing finite number of *k*-grams + tier induction
- learnable in the limit from positive text

Jardine & Heinz (2016)

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Mappings	s we use			

General assumption: we assume stem not to be bound in length:

- There is no limit on the length of the stem in languages.
- The stem can be result of the compounding. whiteboard, whiteboard marker, whiteboard marker cleaning fluid, whiteboard marker cleaning fluid purchase receipt
- Mapping of the stem to a single symbol will result in insensibility to compounds.

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- Mapping of the stem to a single symbol will result in insensibility to compounds.
 - Affixes: affix-to-symbol mapping
 - Stems: symbol-to-symbol mapping

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
000000	0●000	00000000	00000	00
Strictly Lo	ocal Morphology:	affixation		

Example (prefix 'za-', Russian)					
• zaexat'					
'call on the way'					
axxxx					
Bigram * <i>xa</i> ensures that 'za' is a prefix.					

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
000000	0●000	00000000	00000	00
Strictly L	ocal Morphology:	affixation		

Example (prefix 'za-', Russian)					
• exat'	• zaexat'				
'go, drive'	'call on the way'				
XXXX	axxxx				
Bigram $*xa$ ensures that 'za' is a prefix.					

Example (suffix '-s', I	English)			
• dog	• dog <mark>s</mark>			
XXX	xxxb			
Bigram * <i>bx</i> ensures that 's' is a suffix.				

Preliminaries 000000	Lower bound (SL) 00●00	Upper bound (TSL) 00000000	Typology 00000	Conclusion 00
Strictly L	ocal Morphology	: affixation [co	ont.]	
Exampl	e (affixation, English)			
• loc xx		 blacklist xxxxxxx 		
	lockable xxx <mark>b</mark>	• unblack axxxxxx		

Preliminaries 000000	Lower bound (SL) 00●00	Upper bound (TSL) 00000000	Typology 00000	Conclusion
Strictly Lo	ocal Morphology:	affixation [cont.]	
Example	e (affixation, English)			
 lock 	<	 blackl 	ist	
XXX	Х	XXXXX	XXXX	
• unlo	ock <mark>able</mark>	unbla	cklist <mark>able</mark>	
axx	xxb	axxxx	xxxxb	

 $\mathsf{SLG} = \{ \rtimes \mathsf{b}, \ \mathsf{ba}, \ \mathsf{bx}, \ \mathsf{xa}, \ \mathsf{a} \ltimes \}$

This grammar necessarily generates the following forms of English, too: $\forall axxxxx \Leftrightarrow and \forall xxxxxb \ltimes$.

Prelim 0000	ninaries 000	Lower bound (SL) 00●00	Upper bound (TSL) 00000000	Typology ooooo	Conclusion 00
Sti	rictly Loca	al Morphology:	affixation	[cont.]	
	Example (a	affixation, English)			
	 lock 		• bla	cklist	
	XXXX		XXX	XXXXXX	
	• unlock	cable		lacklistable	
	axxxx	0	axx	xxxxxxb	

 $\mathsf{SLG} = \{ \rtimes \mathsf{b}, \ \mathsf{ba}, \ \mathsf{bx}, \ \mathsf{xa}, \ \mathsf{a} \ltimes \}$

This grammar necessarily generates the following forms of English, too: $\forall axxxxx \Leftrightarrow and \forall xxxxxb \Leftrightarrow$.

Indeed, this prediction is correct:

Example (affixation, English)		
• unleash	• breakable	
axxxxx	xxxxxb	



- English un-...-able are prefix and suffix that can co-occur.
- However, two parts of a *circumfix* cannot occur independently:

Consider the following example from Indonesian:

Example (circumfix 'ke-an', Indonesian)			
 tinggi 'high' xxxxxx 	 mahasiswa 'big pupil (student)' xxxxxxxxx 		
 ketinggian 'altitude' axxxxxxb 	 kemahasiswaan 'student affairs' axxxxxxxxb 		
• *axxxxx	• *xxxxxb		

Preliminaries 000000	Lower bound (SL) 0000●	Upper bound (TSL) ೦೦೦೦೦೦೦೦	Typology 00000	Conclusion
SL is not	enough: Indone	esian circumfixa	tion [cont.]	
Example	e (circumfix 'ke-an',	Indonesian)		
• ting xxx	ggi xxxx	• mahasis xxxxxxx		
	inggi <mark>an</mark> xxxxxb	• <mark>ke</mark> maha axxxxxx		
• * <mark>a</mark> >	xxxxx	• *xxxxxx	xb	

 $\begin{aligned} \mathsf{SLG} &= \{ \rtimes \mathsf{b}, \ \mathsf{ba}, \ \mathsf{bx}, \ \mathsf{xa}, \ \mathsf{a} \ltimes \} \\ \mathsf{String} \ \mathsf{language} &= \rtimes \mathsf{xxxx} \ltimes, \ \rtimes \mathsf{axxxxb} \ltimes, \ \rtimes \mathsf{axxxx} \ltimes, \ \rtimes \mathsf{xxb} \ltimes \ldots \end{aligned}$

Preliminaries 000000	Lower bound (SL) 0000●	Upper bound (TSL) 00000000	Typology 00000	Conclusion
SL is not	enough: Indone	sian circumfixat	tion [cont.]	
Example	e (circumfix 'ke-an',	Indonesian)		
• tin; xxx	ggi «xxx	 mahasisy xxxxxxxx 		
	inggi <mark>an</mark> xxxxx <mark>b</mark>	 kemahas axxxxxxx 		
• * <mark>a</mark> >	xxxxxx	• *xxxxxx	xb	

 $\begin{aligned} \mathsf{SLG} &= \{ \rtimes \mathsf{b}, \ \mathsf{ba}, \ \mathsf{bx}, \ \mathsf{xa}, \ \mathsf{a} \ltimes \} \\ \mathsf{String} \ \mathsf{language} &= \rtimes \mathsf{xxxx} \ltimes, \ \rtimes \mathsf{axxxxb} \ltimes, \ \rtimes \mathsf{axxxx} \ltimes, \ \rtimes \mathsf{xxb} \ltimes \ldots \end{aligned}$

Problem:

- SL languages can only capture local dependencies
- Circumfixes introduce non-local ones

Preliminaries 000000	Lower bound (SL) 00000	Upper bound (TSL) ●0000000	Typology 00000	Conclusion
Morphota	actics is TSL			
Exampl	e (circumfix 'ke-an'	, Indonesian)		
• tinggi		• mahasisv xxxxxxxx		
 ketinggian axxxxxxb 			 kemahasiswaan axxxxxxxxxb 	
• *axxxxx		• *xxxxxxx	:b	

 $\mathsf{TSLG}(\mathsf{circumfix_tier}) = \{ \rtimes \mathsf{b}, \mathsf{ ba}, \mathsf{a} \ltimes \}$

Preliminaries 000000	Lower bound (SL) 00000	Upper bound (TSL) ●0000000	Typology 00000	Conclusion 00
Morphota	octics is TSL			
Example	e (circumfix 'ke-an'	, Indonesian)		
• tin	ggi	 mahasiswa 		

 xxxxx
 ketinggian axxxxxb
 *axxxxxb
 *axxxxxb
 *xxxxxxb
 *xxxxxxb

 $\mathsf{TSLG}(\mathsf{circumfix_tier}) = \{ \rtimes \mathsf{b}, \, \mathsf{ba}, \, \mathsf{a} \ltimes \}$

Licit strings: • ×xxxxx • • × • Illicit strings: • ×axxxx • • × axxxx • •

● ⋊*axxxxxb*⋉ ⋊*ab*⋉

● ×bxxxa× ×ba× Preliminaries Lower bound (SL) Upper bound (TSL) Typology Conclusion

Morphotactics is TSL

Example (circumfix 'ka-an', llocano)

In Ilocano, it is impossible to do embedded circumfixation:

 bigát 'morning' xxxxx kabigátan
 'the next morning'
 axxxxxb







 $\mathsf{TSLG}(\mathsf{circumfix_tier}) = \{ \rtimes \mathsf{b}, \mathsf{ba}, \mathsf{a} \ltimes, \mathsf{aa}, \mathsf{bb} \}$





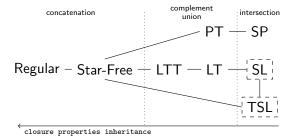
 $\mathsf{TSLG}(\mathsf{circumfix}_{\mathsf{tier}}) = \{ \rtimes \mathsf{b}, \mathsf{ba}, \mathsf{a} \ltimes, \mathsf{aa}, \mathsf{bb} \}$

Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Interim S	ummary			

- SL enforces local dependencies
- TSL enforces local dependencies on the determined tier
- Most of morphotactics is SL, some of it is TSL
- Learning of TSL languages is possible from positive data only
- Can morphotactics be more than TSL?



Can morphotactics be more than TSL?





 Closure under concatenation: Frenglish contains only words whose first part is a word of French and the second a word of English.



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- Closure under concatenation: Frenglish contains only words whose first part is a word of French and the second a word of English. X
- **Closure under union**: If a Mandaresian word violates rules of Mandarin Chinese, it must obey the rules of Indonesian.



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- **Closure under relative complement**: Hsilgne contains all words that are ill-formed in English.



- Closure under concatenation: Frenglish contains only words whose first part is a word of French and the second a word of English. X
- Closure under union: If a Mandaresian word violates rules of Mandarin Chinese, it must obey the rules of Indonesian. X
- Closure under relative complement: Hsilgne contains all words that are ill-formed in English. X



• **Closure under intersection**: Russenorsk is created by combination of elements of Russian and Norwegian.



 ● Closure under intersection: Russenorsk is created by combination of elements of Russian and Norwegian. ✓ (spoken in Northern Norway, 18th-19th centuries)



 Closure under intersection: Russenorsk is created by combination of elements of Russian and Norwegian. (spoken in Northern Norway, 18th-19th centuries)

Example (Closure under intersection)

- A language allows complex nuclei and blocks codas (Supyire)
- A language forbids complex nuclei and allows codas (Russian)



 Closure under intersection: Russenorsk is created by combination of elements of Russian and Norwegian. ✓ (spoken in Northern Norway, 18th-19th centuries)

Example (Closure under intersection)

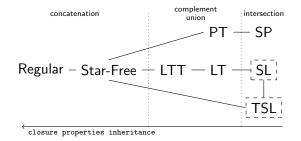
- A language allows complex nuclei and blocks codas (Supyire)
- A language forbids complex nuclei and allows codas (Russian)
- Then there will be a language that blocks complex nuclei and codas (Hawaiian, Senufo)



- X Closure under concatenation
- X Closure under union
- X Closure under relative complement
- Closure under intersection



- X Closure under concatenation
- X Closure under union
- X Closure under relative complement
- Closure under intersection



Preliminaries	Lower bound (SL)	Upper bound (TSL)	Typology	Conclusion
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Typologic	al gaps			

Basic Logic of Argument

- All attested morphotactic patterns must be TSL.
- So if pattern A is TSL, and pattern B is TSL, but their combination A+B is not, we get a typological gap.

Some predicted gaps:

- No embedded circumfixation;
- No cases when amount of affixes A depends on the amount of affixes B;
- In general, no $a^n b^n$ pattern and its derivatives.

Preliminaries 000000	Lower bound (SL) 00000	Upper bound (TSL) 00000000	Typology o●ooo	Conclusion 00
Typologic	cal gap I: Impos	sible compound	ing	
Russia	n pattern – (stem-	o)*-stem		
Exampl	e (compounding, R	ussian)		
	d <mark>o</mark> voz	• vod <mark>o</mark> voz		
'wa	ater carrier'	'carrier	of water carrie	ers'
XXX	XOXXX	XXX <mark>0</mark> XXX	OXXX	

Preliminaries 000000	Lower bound (SL) 00000	Upper bound (TSL) 00000000	Typology o∙ooo	Conclusion
Typological	gap I: Impos	sible compound	ing	
Russian p	oattern – (stem-o)*-stem		
Example (compounding, Ru	issian)		
• vodov		• vodovoz		
'wate	r carrier'	'carrier	of water carriers	,
XXXOX	XX	XXXOXXX	OXXX	

Turkish pattern – stem-(stem⁺-o)

Example (compounding, Turkish)

- bahçe kapı-sı 'garden gate' xxxxxxxxo
- türk kahve-sı 'Turkish coffee'

XXXXXXXXXO

- türk bahçe kapı-sı 'Turkish garden gate' xxxxxxxxxxo
- *türk bahçe kapı-sı-sı *xxxxxxxxxxxx00



Russian pattern – (stem-o)*-stem Turkish pattern – stem-(stem⁺-o)

Turkussian pattern: amount of compound markers is equal to the amount of added stems, stem-(stemⁿ-oⁿ)



Typological gap I: Impossible compounding

Russian pattern – (stem-o)*-stem Turkish pattern – stem-(stem⁺-o)

Turkussian pattern: amount of compound markers is equal to the amount of added stems, $stem-(stem^{n}-o^{n})$

- This pattern is not regular because it has infinite number of "good continuations". (*Myhill-Nerode theorem*)
- It appears to be non-existent.



Typological gap II: Recurrent affixation

Sometimes languages allow some affixes to be iterated: a*-stem.

Consider example of such pattern in German:

Example (prefix 'über', Gern	nan)
 morgen 'tomorrow' 	 übermorgen 'the day after tomorrow'
• überübermorgen 'the day after the day a aaxxxxxx	axxxxx



German pattern: a*-stem.

The same meaning can be expressed in another language differently, consider llocano (Austronesian) temporal circumfix *ka-...-an* 'next'.

Example (circumfix 'ka-a	n', llocano)
 bigát 	• kabigátan
'morning'	'the next morning'
XXXXX	axxxxp



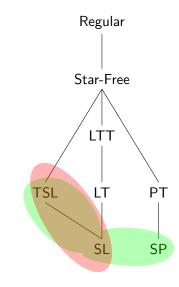
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The same meaning can be expressed in another language differently, consider llocano (Austronesian) temporal circumfix *ka-...-an* 'next'.

Example (circumfix 'ka-a	n', llocano)
• bigát	 kabigátan
'morning'	'the next morning'
XXXXX	axxxxb

However, word kakabigátanan doesn't appear to be possible word in llocano: aⁿ-stem-bⁿ pattern is not regular.





- Morphotactics is at most Tier-Based Strictly Local
- Positive data is enough for morphological learning
- Set of typological gaps can be explained due to the subregular nature of morphology
- Same formal tools can be used for morphology and phonology

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Future wor				

- Try to find SP patterns in morphotactics
- Look at more typologically diverse languages
- Extend to mappings from underlying to surface forms
- Work with representations of internal structure
- The elephant in the room: reduplication

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Thank you!

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