Concl

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**RS** Constraints

- Reference-Set Constraints
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### Reference-Set Constraints

Reference-set constraint  $\approx$  transderivational constraint  $\approx$  global economy condition  $\approx$  interface strategy

#### An Informal Definitior

Given some input tree t, a reference-set constraint computes a set of possible output trees for t — called the reference set of t — and picks from said set the optimal output tree according to some economy metric.

Some examples from the literature:

- Rule I (Reinhart 2006)
- Scope Economy (Fox 2000)
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## Example: Focus Economy

- (1)a.  $[_{TP} \text{ John } [_{VP} \text{ bought } [_{DP} \text{ a red } \mathbf{car}]]].$ Focus set: {TP, VP, DP, red car, car}
  - b.  $[_{TP}$  John  $[_{VP}$  bought  $[_{DP}$  a **red** car]]]. Focus set: {red}

Any constituent containing the carrier of sentential main stress may be focused.

If the main stress has been shifted, a constituent containing its carrier may be focused iff it cannot be focused in the tree with unshifted stress.

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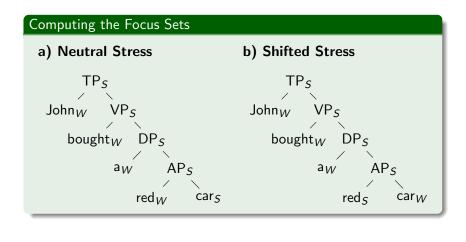
- (2) a. [TP John [VP bought [DP a red car]]]. Focus set: {TP, VP, DP, red car, car}
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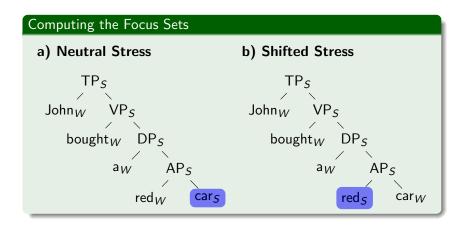
### Focus Projection

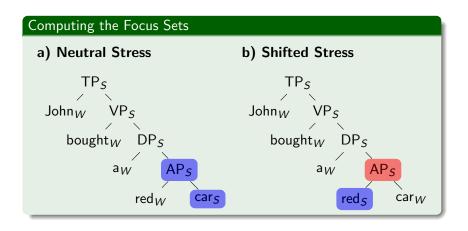
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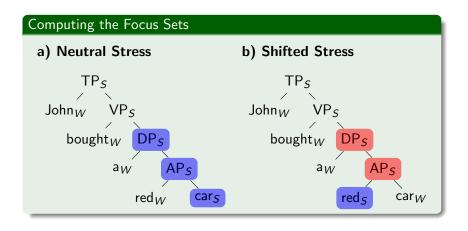
### Focus Economy Rule

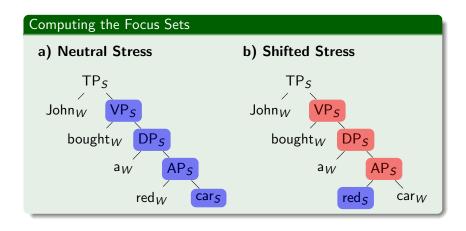
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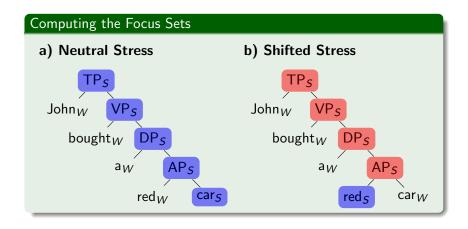












RS Constraints

## • Reference-set constraints are argued to be

- too computationally demanding for narrow syntax
- (Collins 1996; Johnson and Lappin 1999)
- unwieldy for empirical work (Sternefeld 1996)
- But if we use linear tree transducers as a model, it turns out that the concerns are unwarranted. Rather, reference-set constraints are...
  - Unity 1 (theory-internal)

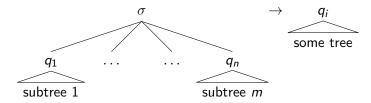
     a different way of specifying standard well-formedness
     constraints
  - Unity 2 (across theories)
     connected to OT and Synchronous TAG.
  - Unity 3 (connection of formal and empirical work) ideally suited to account for cross-linguistic variation (which they are hardly ever used for).

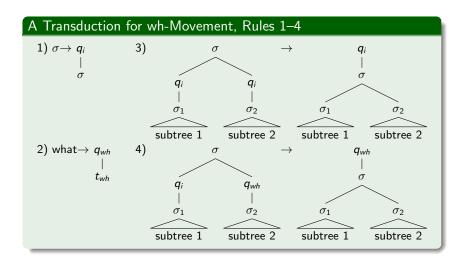
### Linear Tree Transducers in Pictures

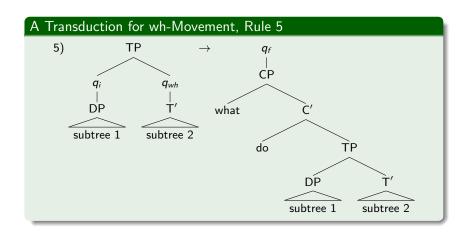
### A linear finite-state bottom-up tree transducer

- traverses an input-tree from the leaves towards the root,
- labels it with states  $q_i$ , and
- transforms it into an output-tree.

It does so using rules of the following kind:







# A Transduction for wh-Movement, Application TP DP the men what like

## A Transduction for wh-Movement, Application TP DP the men like q<sub>wh</sub> $t_{wh}$

## A Transduction for wh-Movement, Application TP DP the men $q_i$ q<sub>wh</sub> like $t_{wh}$

# A Transduction for wh-Movement, Application TP DP the men $q_{wh}$ like

## A Transduction for wh-Movement, Application TP DP the men $q_{wh}$ **VP** like

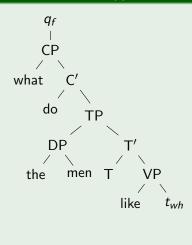
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### A Transduction for wh-Movement, Application



## Some Important Facts

### What is Possible?

- Relabeling nodes
- Deleting subtrees
- Inserting subtrees of bounded size
- Enforcing constraints that define regular tree languages

### What is Impossible?

- Copying of arbitrary subtrees
- Switching positions of non-siblings (e.g. specifier and complement)
- Counting past some threshold

### Mathematical Properties

- A transducer can be decomposed into a sequence of smaller transducers, et vice versa.
- If the input tree language of a transducer is regular, then so is its output language. Regular tree languages are sufficiently powerful for Minimalism (Kobele et al. 2007).

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## Overall Reasoning

### Strategy

For a given reference-set constraint C, exhibit

- a Minimalist grammar that generates the input language, and
- a sequence of transducers that computes the same mapping from inputs to optimal outputs.
- Due to the mathematical properties of transducers, the output language is no more complex than the input language
- Hence it can be generated by some Minimalist grammar
- Hence C is equivalent to some "constraint" that does not involve reference-set computation.

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But why should this work for arbitrary reference-set constraints?

## OT: A Bird's Eye Perspective

It seems natural to model reference-set constraints via OT.

### Reference-Set Constraints as OT Grammars

- Use GEN to compute the reference-sets.
- Use a sequence of constraints to filter out the suboptimal candidates.

### A Major Complaint

Without further restrictions, OT grammars can generate any kind of (tree) language

⇒ they don't tell us anything about reference-set constraints.

Fortunately, there is a weaker alternative...

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## Optimality Systems: A Restricted Version of OT

### Optimality Systems (OSs; Frank and Satta 1998)

A variant of OT that keeps just the bare skeleton.

- All constraints only consider the output, never the input.
- No correspondence theory
- No output-output correspondence
- No sympathy constraints

There are mathematical conditions that ensure that an OS can be implemented by a tree transducer.

#### A Minor Quibble

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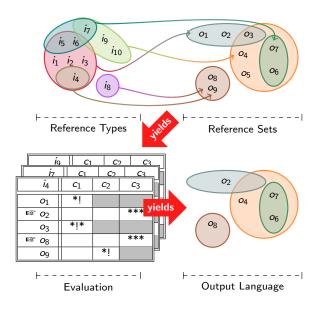
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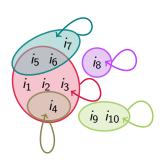
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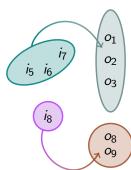
# Depiction of a Controlled OS



### Reference-set Constraints as Controlled OSs

- Almost all constraints in the literature exhibit one of the two configurations below.
- What do the two have in common?





#### Output Joint Preservation

If two reference sets overlap, then so do the reference types that are mapped to them.

#### Theorem (Frank and Satta 1998; Wartena 2000; Jäger 2002)

A controlled OS can be implemented as a transducer if

- the OS is output-joint preserving, and
  - the input language is regular, and
- GEN and all constraints can be implemented as transducers

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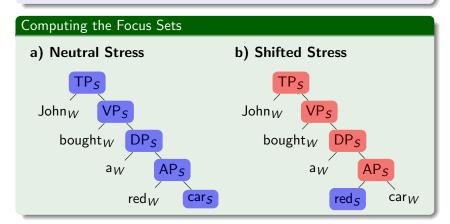
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# Example 1: Focus Economy

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If the main stress has been shifted, a constituent containing its carrier may be focused iff it cannot be focused in the tree with unshifted stress.

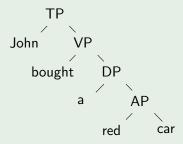


### Transducer Model: GEN

#### Step 1 & 2: GEN

- Non-deterministically relabel input with S/W-subscripts.
- Non-deterministically focus some node along the "stress path".

### Transducing an Input into a Stress-Annotated Output with Focus

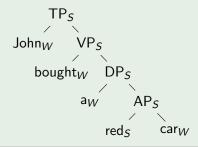


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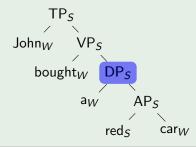


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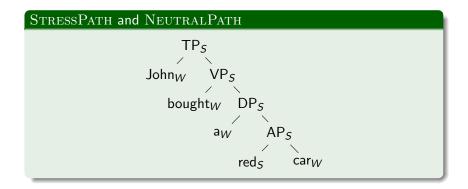
### Transducer Model: The Constraint

Focus Economy requires reference to the neutral stress pattern. We allow this by implicitly representing the neutral stress within the same tree!

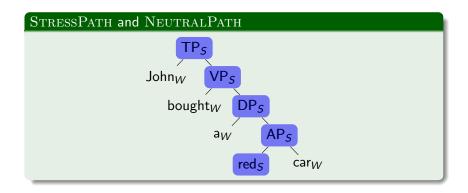
### Strategy

- Define two paths STRESSPATH and NEUTRALPATH.
- StressPath represents the path of the current stress.
- $\bullet$   $\ensuremath{\mathrm{NEUTRALPATH}}$  represents the path of the neutral stress.
- Add a constraint that requires focus to be in the stress path, but unless StressPath and NeutralPath pick out the same nodes, focus may not be in NeutralPath.

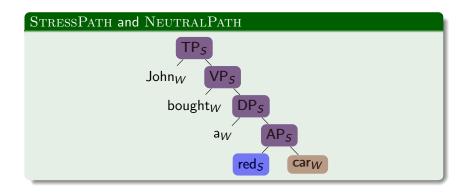
# Example of $\phi$



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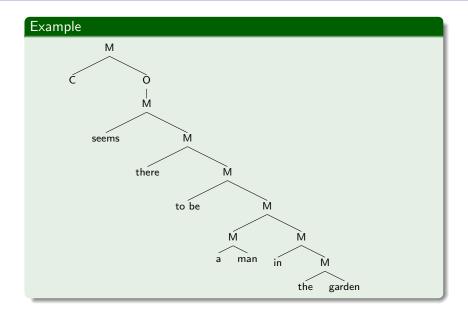
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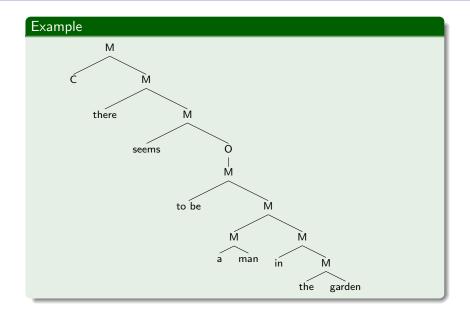
If two convergent derivations d and d' are built from the same lexical items and identical up to step n, at which point d continues with Merge and d' with Move, filter out d'.

- (3) a. There seems  $t_{\text{there}}$  to be a man in the garden.
  - b. \* There seems a man to be  $t_{a \text{ man}}$  in the garden.
  - c. A man seems  $t_{a \text{ man}}$  to be  $t_{a \text{ man}}$  in the garden.

# Derivation Trees of (3a) and (3b)

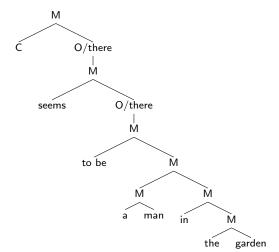


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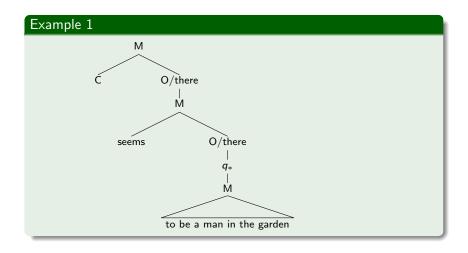
# Transducer Model: GEN (Step 1)

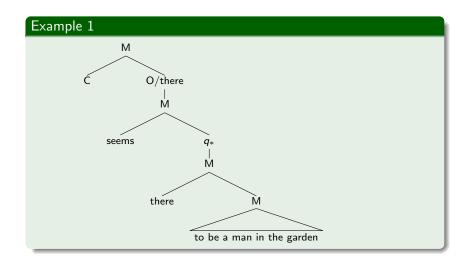
- Fuse the two derivations into one underspecified derivation.
  - Remove all features but the category feature.
  - Inside TP: Replace O or Merger of there by new label O/there.

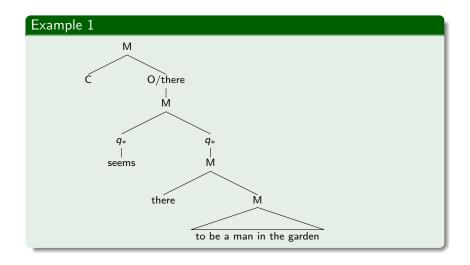


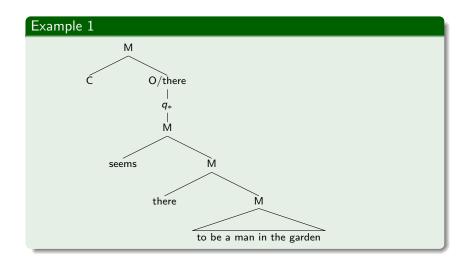
# Transducer Model: GEN (Step 2)

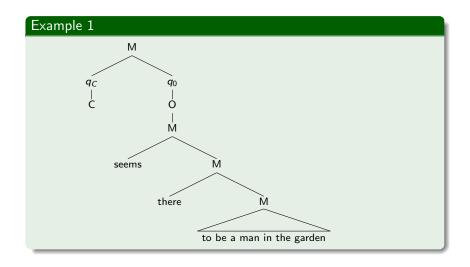
- Turn O/there back into O or Merge of there.
  - Use a transducer with states  $q_*$ ,  $q_O$  and  $q_C$ .
  - In state q<sub>\*</sub>, the transducer non-deterministically rewrites
     O/there as O or Merge of there.
  - If the transducer rewrites O/there as O, it switches into state q<sub>0</sub>.
  - In state  $q_0$ , every occurrence of O/there is rewritten just as O.
  - The transducer switches out of  $q_0$  only if it encounters a CP (indicated by state  $q_C$ ; cf. structured numerations).
- Reinstantiate the features.

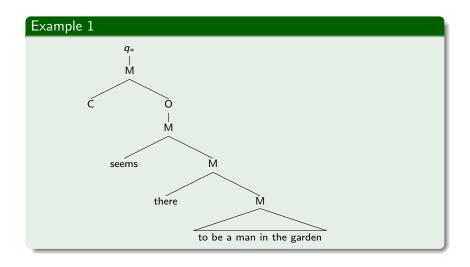


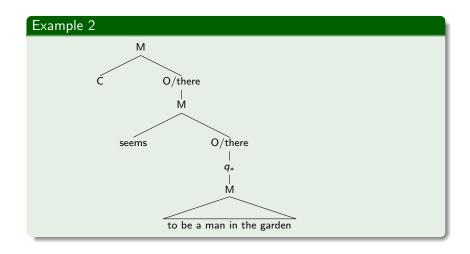


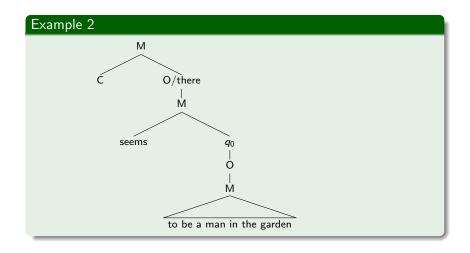


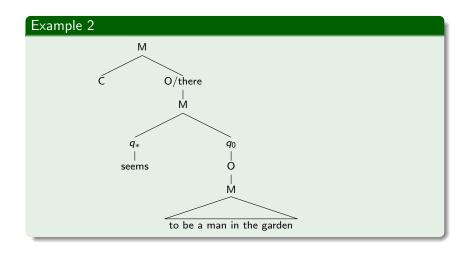


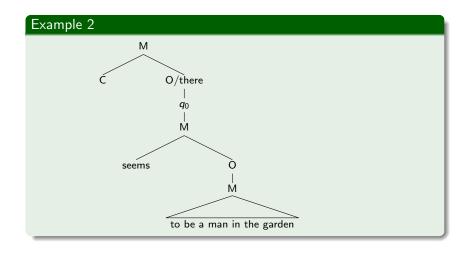


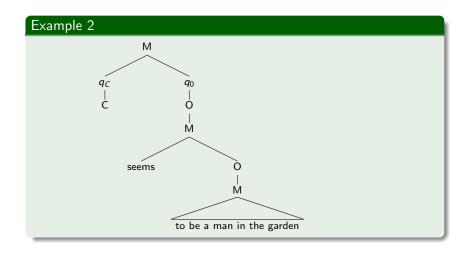


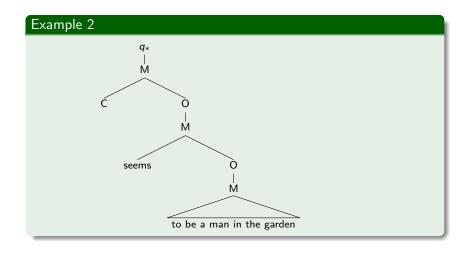












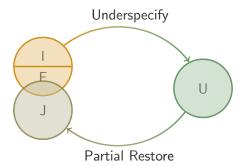
# Transducer Model: The Induced Mapping

The output candidates for both (4a) and (4b) are now (5a)–(5b).

- (4) a. There seems  $t_{\text{there}}$  to be a man in the garden.
  - b. \* There seems a man to be  $t_{a \text{ man}}$  in the garden.
- (5) a. \* There seems there to be a man in the garden.
  - b. There seems  $t_{\text{there}}$  to be a man in the garden.
  - c. A man seems  $t_{a \text{ man}}$  to be  $t_{a \text{ man}}$  in the garden.
- We may extend the mapping such that (5c) is also assigned this reference set.
- (5a) still has to be ruled out.

### Transducer Model: The Constraint

The only constraint is the input language itself! By turning it into a transducer and composing it with GEN, we remove all instances of overgeneration and filter out the illicit MOM violators.



# Shortest Derivation Principle (SDP)

#### **SDP**

Given convergent derivations  $d_1, \ldots, d_n$  over the same lexical items, pick the one(s) with the fewest instances of Move.

Why is the following sentence ungrammatical?

(6) \* Who<sub>i</sub> was [ $DP_i$  a picture of  $t_i$ ] taken  $t_i$  by John?

# Derivations for (6)

RS Constraints

Two derivations are possible for (6).

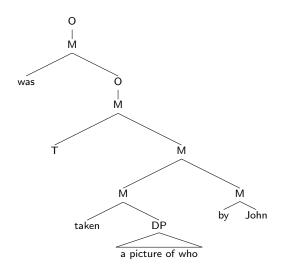
CED violation in (7c)

- (7) a.  $[VP \text{ taken } [DP_i \text{ a picture of who}_i]$  by John]
  - b.  $[TP [DP_i]$  a picture of who<sub>i</sub> ]T [VP] taken  $t_i$  by John]
  - c. [CP] who; was [TP] [DP] a picture of  $t_i$  ] T [VP] taken  $t_i$ by John]]]

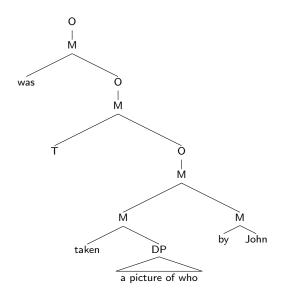
Derivation (8) is longer than (7)!

- a. [VP] taken [DP] a picture of who; by John (8)
  - b. [VP] who; taken [DP]; a picture of  $t_i$ ] by John]
  - c.  $[TP]_{DP_i}$  a picture of  $t_i$   $T[_{VP}$  who<sub>i</sub> taken  $t_i$  by John]
  - d. [CP who; was [TP [DP; a picture of  $t_i$ ] T [VP taken  $t_i$ by John]]]

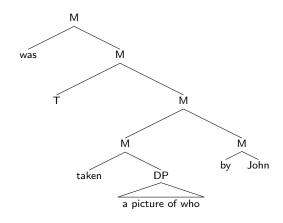
# Derivation Tree of (7)



# Derivation Tree of (8)



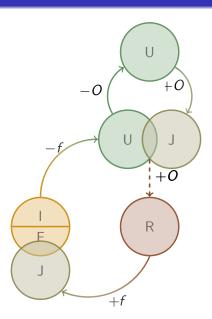
## Underspecified Derivation Tree of (7) and (8)



## Strategy

- Compute reference-set by
  - mapping to underspecified derivation (i.e. remove Move-nodes)
  - arbitrarily adding Move-nodes to underspecified derivation
  - discarding all derivation trees that aren't in the input language (i.e. the junk)
- Filter out the suboptimal derivation trees (those that can be obtained from others by adding Move-nodes)
  - Let R be the transduction that maps a derivation tree to the trees in its reference-set and +O the transduction defined by adding Move-nodes
  - ② The range of the composition of R with +O is the set of derivation trees that can be obtained from some tree in the range of R by adding Move-nodes, i.e. the suboptimal outputs.
  - 3 Thus, the relative complement of the range of R and the range of the composition of R with +O is the set S of optimal outputs. Composing R with the diagonal over S maps every tree to its optimal outputs.

### Architecture of SDP



## Underspecification-and-Filtration

#### A Rule of Thumb

A reference-set constraint is likely to be computable by

- a transducer if
  - one can find a structure that encodes the commonalities of all the competitors, and
  - neither the underspecification step nor the recovery step require insertion of material of unbounded size, and
  - the economy metric can be implemented as
    - a well-formedness constraint on underspecified structures, or
    - a specific restriction on the recovery step, or
    - a transducer that turns optimal candidates into suboptimal ones.

RS Constraints Transducers General Results Focus Economy MOM SDP Concl References

# Why Keep Them?

### Modularity

Constraint only depends on input language, not on mechanisms that generate it

#### Succinctness

Non-reference-set correspondent may fail to make the restriction explicit or be much more complicated; reference-set constraint may subsume very different constraints

#### More Tweakable Parameters

Reference-set constraint gives us at least four parametrizable components: reference types, reference sets, the map between the two, and the economy metric.

### Reaching out

Connections to OT, sTAG and others may allow us to incorporate results from these frameworks

# Conclusion (Part 1)

- Tree transducers were proposed as a model for reference-set constraints.
- OSs offer a bird's eye view on them (Unity 2).
- Most requirements for an OS to be efficiently computable are fulfilled by reference-set constraints; in particular, their corresponding OSs are output joint preserving.
- ullet The only problematic areas are  $\operatorname{GEN}$  and the OS constraints.
- The underspecification-and-filtration strategy offers a general solution.

# Conclusion (Part 2)

### Concern 1: Computability

- If a reference-set constraint can be modeled as a tree transducer, it is efficiently computable.
- In fact, it is equivalent to some standard well-formedness condition (**Unity 1**).

### Concern 2: Empirical applicability

- Now that we have a natural class of licit reference-set algorithms and economy metrics (those definable by tree transducers), we can see if more nuanced revisions of them are more suitable.
- We can use computational tools to test our predictions.
- The additional parameters may make available new cross-linguistic generalizations (Unity 3).

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