Subregular Morpho-Semantics The Expressive Limits of Monomorphemic Quantifiers

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Stony Brook University mail@thomasgraf.net http://thomasgraf.net Rutgers University December 15, 2017 Vou can get the slides here under "News" Take-Home Message

- Supplement linguistic theory with computational perspective
- Typological gaps can be explained computationally.

Case Study: Morphosemantics of Quantifiers

A D-quantifier may have a monomorphemic realization only if its quantifier language is TSL.

Outline

1 TSL Patterns in Phonology, Morphology, and Syntax

2 TSL Morpho-Semantics

- Quantifier Languages
- All Monomorphemic Quantifiers are TSL
- Tightening the Characterization

3 A Broader Program

Subregular Hierarchy

- 1 define different classes of grammars
- 2 organize these classes into an expressivity hierarchy
- 3 needed level of expressivity?

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- 3 needed level of expressivity?



TSL: Tier-Based Strictly Local

- All patterns described by markedness constraints that are
 - inviolable,
 - locally bounded,
 - formalized as n-grams.
- Non-local dependencies are local over tiers. (Goldsmith 1976)
- Linguistic core idea:

Dependencies are local over the right structure.

- Captured by forbidding voiced segments at the end of a word
- German: Don't have z\$ or v\$ or d\$ (where \$ = word edge).

Example: German		
	*z\$	
* \$ r a d \$	* v\$	
	* d\$	

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- Captured by forbidding voiceless segments between vowels
- Suppose:
 - ► $[-\text{voice}] = \{s, \int\}$
 - ► $V = {a,i,u}$
- ► Then: don't have asa, a∫a, asi, a∫i, ...

Example

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TSL

Adding Tiers: Samala Sibilant Harmony

- ► If multiple sibilants occur in the same word, they must all be +anterior (s,z) or -anterior (∫,3).
- In other words: Don't mix purple and teal.

But: Sibilants can be arbitrarily far away from each other!

Example: Samala

```
*$hasxintilawa∫$
```

```
$ha∫xintilawa∫$
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*\$<mark>s</mark>tajanowonwa∫\$

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Example: Samala

- Let's take a hint from phonology: create locality with a tier. (Heinz et al. 2011)
- Restriction 1: only 1 tier
- Restriction 2: projection is determined by the segments, not their environment



Jeff Heinz

Example: Samala Revisited

```
Project sibilant tier
*sſ, *sʒ, *zʃ, *zʒ, *ʃs, *ʒs, *ʃz, *ʒz
```

*\$hasxintilawa\$\$ \$ha\$xintilawa\$\$

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Example: Samala Revisited

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    Project sibilant tier
```

2 *s∫, *sȝ, *z∫, *zȝ, *∫s, *ȝs, *∫z, *ȝz

*\$ha<mark>s</mark>xintilawa∫\$\$ha∫xintilawa∫\$

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$ $ $

$ l
|

*$ hasxintilawaʃ$
$ haſxintilawaʃ$
```

TSL

Culminativity: Simple Counting with TSL











Why is TSL Interesting?

- Linguistically natural
- Correct and very efficient learning algorithm (Jardine and McMullin 2017)
- Low resource demands \Rightarrow cognitively plausible
- Captures wide range of phonotactic dependencies
- Cannot generate many unattested patterns

- Harmony only holds between initial and final segments
- Linguistically plausible, yet unattested

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TSL Semantics

- TSL seems to play an important role in
 - phonology,
 - morphology,
 - syntax.
- What's missing? Semantics!
- But TSL is about strings/trees.
- What is a semantic string language?

Formal Language Theory for Semantics

Quantifier Languages Meanings as strings of truth values (van Benthem 1986)

- "Tense Languages" Meanings as strings of events (Fernando 2011)
- I'll only talk about quantifier languages here.
- Ongoing work with Rob Pasternak on subregularity of tense languages



Evaluating the Truth of Quantifiers

- (1) a. Every student cheated.
 - b. No student cheated.
 - c. Some student cheated.
 - d. Three students cheated.

students	John	Mary	Sue
cheated	yes	no	yes
string	Y	Ν	Y

- (1a): False, because the string contains a N
- ▶ (1b): False, because the string contains a Y
- ▶ (1c): **True**, because the string contains a Y
- ▶ (1d): **False**, because the string does not contain three Ys

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Formalization Step 1: Binary String Languages

Idea: Convert relation between sets A and B into set of Yes/No-strings

Definition (Binary String Language)

1 A, B: arbitrary sets

- **2** f(A, B): maps each $a \in A$ to Y if $a \in B$, otherwise N
- **3** e(A): arbitrary enumeration of A
- **4** L(A, B): all e(A), relabeled by f(A, B)

Example

Set of Set of				Mary, Su Sue, Bill,	-
2 f(A , B)): N	lohn ⊢ ⁄Iary ⊢ Sue ⊢	→ N		
3 e(A):	2) 3) 4) 5)	John John Mary Mary Sue Sue	Sue John Sue John	Mary Sue John Mary	
4 L(A,B): {	YNY, YYN, NYY	}		

Formalization Step 2: Quantifier Language

Idea: Every quantifier is a set of acceptable Yes/No-strings

Definition (Quantifier Language)

L(Q) is the **quantifier language** of Q iff it holds for all A and B that Q(A, B) is true iff $L(A, B) \subseteq L(Q)$.

Example

- L(every) = set of all strings containing no N
- ► Why?
 - every(A, B) iff $A \subseteq B$
 - If $A \subseteq B$, then no binary string contains N.
 - If some binary string contains N, then $A \not\subseteq B$.

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Quantifier Constraint every no some at least n at most n exactly n not all all but n most an even number

Quantifier Constraint every $|\mathbf{N}| = 0$ no some at least n at most n exactly n not all all but n most an even number

Quantifier Constraint every $|\mathbf{N}| = 0$ $|\mathbf{Y}| = 0$ no some at least n at most n exactly n not all all but n most an even number

an

Quantifier	Constraint
every	$ \mathbf{N} = 0$
no	$ \mathbf{Y} = 0$
some	$ \mathbf{Y} \ge 1$
at least <mark>n</mark>	
at most <mark>n</mark>	
exactly <mark>n</mark>	
not all	
all but <mark>n</mark>	
most	
even number	

Quantifier	Constraint
every	$ \mathbf{N} = 0$
no	$ \mathbf{Y} = 0$
some	$ \mathbf{Y} \ge 1$
at least <mark>n</mark>	$ \mathbf{Y} \ge \mathbf{n}$
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not all	
all but <mark>n</mark>	
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at most <mark>n</mark>	$ \mathbf{Y} \leq \mathbf{n}$
exactly <mark>n</mark>	
not all	
all but <mark>n</mark>	
most	
an even number	

Quantifier	Constraint
every	$ \mathbf{N} = 0$
no	$ \mathbf{Y} = 0$
some	$ \mathbf{Y} \ge 1$
at least <mark>n</mark>	$ \mathbf{Y} \ge \mathbf{n}$
at most <mark>n</mark>	$ \mathbf{Y} \leq \mathbf{n}$
exactly n	$ \mathbf{Y} = \mathbf{n}$
not all	
all but <mark>n</mark>	
most	
an even number	

Quantifier	Constraint
every	$ \mathbf{N} = 0$
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all but <mark>n</mark>	
most	
an even number	

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not all	$ \mathbf{N} \ge 1$
all but <mark>n</mark>	$ \mathbf{N} = \mathbf{n}$
most	
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at most <mark>n</mark>	$ \mathbf{Y} \leq \mathbf{n}$
exactly n	$ \mathbf{Y} = \mathbf{n}$
not all	$ \mathbf{N} \geq 1$
all but <mark>n</mark>	$ \mathbf{N} = \mathbf{n}$
most	$ \mathbf{Y} > \mathbf{N} $
an even number	

Constraint
$ \mathbf{N} = 0$
$ \mathbf{Y} = 0$
$ \mathbf{Y} \ge 1$
$ \mathbf{Y} \ge \mathbf{n}$
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$ \mathbf{Y} = \mathbf{n}$
$ \mathbf{N} \geq 1$
N = n
$ \mathbf{Y} > \mathbf{N} $
$ \mathbf{Y} $ even

TSL Quantifier Languages for every and no



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\$ N N N N \$ \$ N N Y N \$



***Y**

\$ N N N N \$

\$ N N Y N \$

























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SL



SL

Most Quantifier Languages Require a Tier



*\$\$, *\$N\$, *\$NN\$, *NNNN

\$ N N Y N \$ \$ N Y Y N \$ \$ N N N N N \$



iL



5L



\$NN N\$	\$N N \$	\$ N N N N \$
\$ N N Y N \$	\$ N Y Y N \$	\$ N N N N \$



*\$\$, *\$N\$, *\$NN\$, *NNNN

\$ N N N \$	\$ N N \$	\$ N N N N \$
\$ N N Y N \$	\$ N Y Y N \$	\$ N N N N \$



\$ <mark>NN N\$</mark>	\$N N \$	\$ N N N N \$
\$ N N Y N \$	\$ N Y Y N \$	\$ N N N N \$

SNNYNS

Most Quantifier Languages Require a Tier



SNYYN

SNNNN



L

Most Quantifier Languages Require a Tier



\$ N N N \$	\$ N N \$	\$ N N N N \$
\$ N N Y N \$	\$ N Y Y N \$	\$ N N N N \$

\$

Quantifier	Constraint	<i>n</i> -grams	Tier	
every	$ \mathbf{N} = 0$	* N	none	
no	$ \mathbf{Y} = 0$	* Y	none	
some	$ \mathbf{Y} \ge 1$	*\$\$	Υ	
at least <mark>n</mark>	$ \mathbf{Y} \ge \mathbf{n}$	*\$ Y ^m \$ (m < n)	Υ	
at most <mark>n</mark>	$ \mathbf{Y} \leq \mathbf{n}$	* Y ⁿ⁺¹	Υ	
exactly <mark>n</mark>	$ \mathbf{Y} = \mathbf{n}$	at least $+$ at most	Υ	
not all	$ N \ge 1$	*\$\$	Ν	
all but <mark>n</mark>	$ \mathbf{N} = \mathbf{n}$	at least $+$ at most	Ν	
Quantifier	TSL?	Tier	Mono.	(Paperno 2011)
-----------------	------	------	-------	----------------
every	yes	none	yes	
no	yes	none	yes	
some	yes	Y	yes	
(at least) two	yes	Y	yes	
(at most) two	yes	Y	yes	
not all	yes	Ν	no	
all but one	yes	Ν	no	
even number	no		no	
prime number	no		no	
infinitely many	no		no	
most	no		???	

Quantifier	TSL?	Tier	Mono.	(Paperno 2011)
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no	yes	none	yes	
some	yes	Υ	yes	
(at least) two	yes	Υ	yes	
(at most) two	yes	Υ	yes	
not all	yes	Ν	no	
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even number	no		no	
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infinitely many	no		no	
most	no		???	

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some	yes	Υ	yes	
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all but one	yes	Ν	no	
even number	no		no	
prime number	no		no	
infinitely many	no		no	
most	no		???	

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The Case of most

There is good semantic evidence that "most" is

internally complex and hence not monomorphemic. (Hackl 2009)

	Quantifier	TSL?	Tier	Mono.
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	no	yes	none	yes
	some	yes	Υ	yes
	(at least) two	yes	Y	yes
	(at most) two	yes	Y	yes
	not all	yes	Ν	no
	all but one	yes	Ν	no
	even number	no		no
	prime number	no		no
	infinitely many	no		no
most		no		no

A New Upper Bound on Typological Variation

TSL Interpretation Conjecture

If a language uses a quantifier as a monomorphemic determiner, then its quantifier language must be TSL.

TSL is Too Large

- ► All monomorphemic quantifiers are TSL.
- But not all TSL-definable quantifiers are monomorphemic.
- Why might that be?

Quantifier	TSL?	Tier	Mono.
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no	yes	none	yes
some	yes	Y	yes
(at least) two	yes	Y	yes
(at most) two	yes	Y	yes
not all	yes	Ν	no
all but one	yes	Ν	no

Definition (Monotonicity)

- ▶ Let **A** and **B** be two sets with orders \leq_A and \leq_B , respectively.
- A function f from A to B is monotonic iff

$$x \leq_{\mathbf{A}} y \Rightarrow \mathbf{f}(x) \leq_{\mathbf{B}} \mathbf{f}(y)$$

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Monotonicity in Language

Monotonicity in phonology

- No Crossing Branches constraint
- Natural classes are convex

Monotonicity in morphology

*ABA

Monotonicity in syntax

- Subcategorization < A-Move < A'-Move</p>
- Adjunct Island Constraint & Coordinate Structure Constraint

Monotonicity in semantics

Everywhere...

- Suppose, then, that monotonicity is a desirable trait.
- How does monotonicity relate to tier projection?



Monotonicity forbids projecting only N.

- Suppose, then, that monotonicity is a desirable trait.
- How does monotonicity relate to tier projection?



Project: Y

- Suppose, then, that monotonicity is a desirable trait.
- How does monotonicity relate to tier projection?



Project: Yand N

- Suppose, then, that monotonicity is a desirable trait.
- How does monotonicity relate to tier projection?



Project: nothing

- Suppose, then, that monotonicity is a desirable trait.
- How does monotonicity relate to tier projection?



Project: forbidden

- Suppose, then, that monotonicity is a desirable trait.
- How does monotonicity relate to tier projection?



Project: forbidden

Adding Tiers for every and no

- every and no are the only quantifiers without tier
- But: no tier = tier containing everything



- ► So *every* and *no* can be viewed as using the tier {Y,N}.
- This satisfies monotonicity.



Remaining Issues & Extensions

- TSL also allows for some unnatural quantifiers; ruling them out requires some stipulations.
- What about fuzzy quantifiers? many, few, ...
- TSL makes cognitive complexity predictions; we're working on experiments.
- Where else in semantics does TSL matter?
 - adverbial quantifiers
 - temporal semantics
 - modals
- But those are just small pieces of a much larger puzzle...

The Bigger Goal

- Computational approaches are abstract and content-neutral.
- This isn't a problem but a virtue.
- Abstraction makes it possible to identify parallels between very different domains.

A Program of Subregular Unification

- To what extent can very different properties of language be reduced to the same computational property?
- What are the implications for
 - typological variation,
 - ▶ learnability,
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Place of Morphosemantics



Conclusion

- Among determiners, all monomorphemic quantifiers have quantifier languages that are TSL.
- The opposite does not hold, additional restrictions on TSL are needed.
- ▶ Why does it matter? Because TSL is everywhere in language.
- Ultimate goal:

computational explanation of typological variation

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